

Simulations distribuées, multi-robots : pourquoi et comment garantir la répétabilité des résultats ?

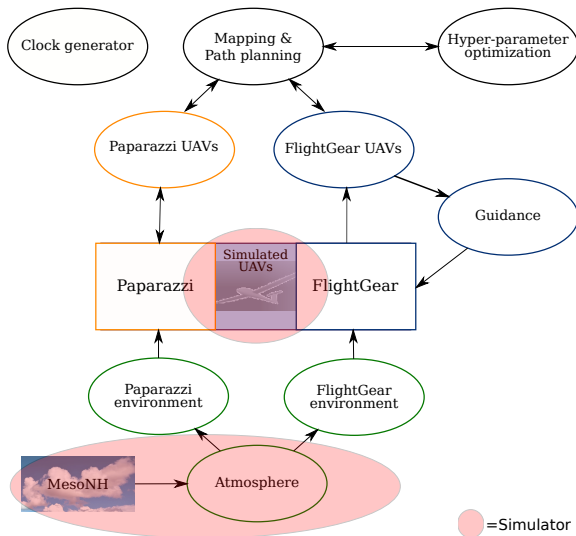
Distributed, multi-robot simulations: why and how to
ensure repeatability of results ?

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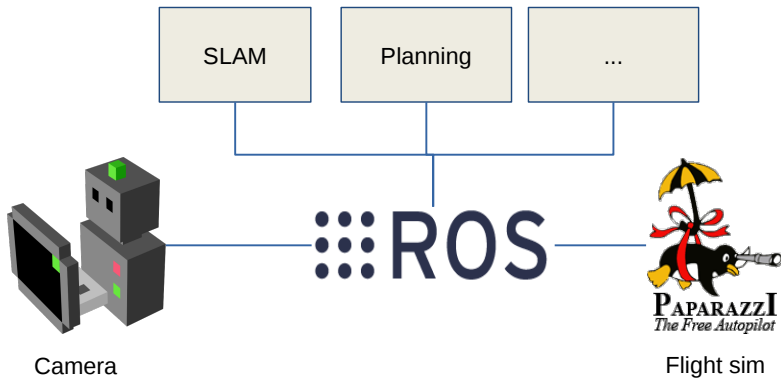
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SKYSCANNER



PRECIDRONE



WHAT DO I WANT FROM THE SIMULATION FRAMEWORK?

- ▶ Compose distributed, modular simulations, interconnect simulators
- ▶ Switch and mix betw. simulated and real components
- ▶ Synchronize simulations, control time advancement
- ▶ Have repeatable simulation results
- ▶ Run realtime simulations
- ▶ Easy to integrate

OVERVIEW

Case studies

Why do we need repeatability?

Why not HLA ?

Proposal: the DSAAM library

Discussion

WHAT IS REPEATABILITY?

The same initial conditions should always produce the same simulation results.

Replicability \neq Reproducibility

Replicability is not Reproducibility: Nor is it Good Science,

C. Drummond, 2009

- ▶ Reproducibility implies introduction of variability
- ▶ Replicability can't be used to prove scientific results
- ▶ Only useful to detect scientific fraud (debatable)

WHAT IS REPEATABILITY'S PURPOSE FOR THE SCIENTIST ?

We are not (only) scientists, but also (mostly) software developpers.

Bug finding and checking:

- ▶ Removes unwanted variability
- ▶ Replicate bugs
- ▶ Regression testing

Ease of development and use:

- ▶ Removes unwanted variability
- ▶ Prototyping with non-realtime algorithms
- ▶ Launching of simulations in batches

Verifiable results ?

IS MY SIMULATION REPEATABLE ?

Probably not, if you are using a robotic simulator.

WHAT DETERS REPEATABILITY ?

System variability:

- ▶ Networking delays
- ▶ Dropped messages (full buffers)
- ▶ System power & load
- ▶ Scheduling policies
- ▶ ...

Interaction simulation \leftrightarrow tested algorithm(s):

- ▶ Out of order message processing
- ▶ Non-deterministic algorithm
- ▶ ...

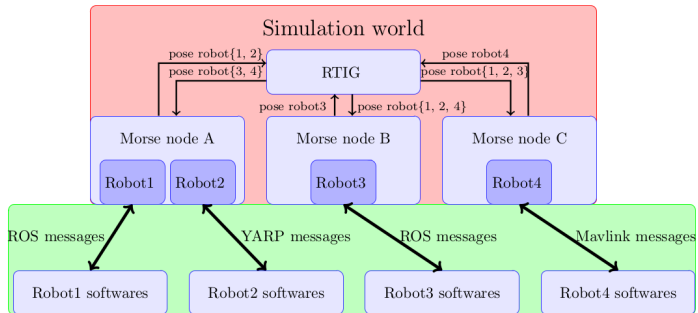
BRIEF INTRO TO HLA

Industry/Military standard

- ▶ Interconnect simulators (exchange state)
- ▶ Manage simulation
- ▶ Time management
- ▶ Best effort (realtime) mode

... what about DIS ?

HLA INTEGRATION IN MORSE



TIME MANAGEMENT IN HLA

- ▶ Sim. wants to advance time to T
- ▶ Sim. sends request
- ▶ RTI delivers all messages up to T
- ▶ Sim. request is granted
- ▶ Sim. advances time to T

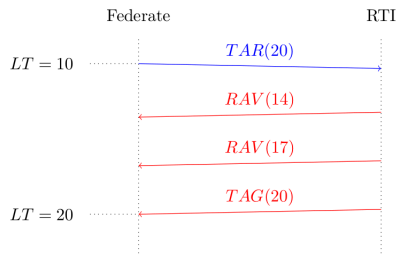


Fig. 2. Typical execution sequence

ADVANTAGES AND DRAWBACKS

Pros:

- ▶ It just works™
- ▶ Provides time management
- ▶ It is an industry standard
- ▶ Centralized architecture

Cons:

- ▶ Military standard
- ▶ Quite complex to use, not many training resources
- ▶ Dependency on the RTI: no API/wire compatibility
- ▶ Adds another middleware for the roboticist
- ▶ No tools, no open-source ecosystem
- ▶ Centralized architecture

DSAAM = Decentralized Synchronisation Architecture for Asynchronous Middleware (name may change)

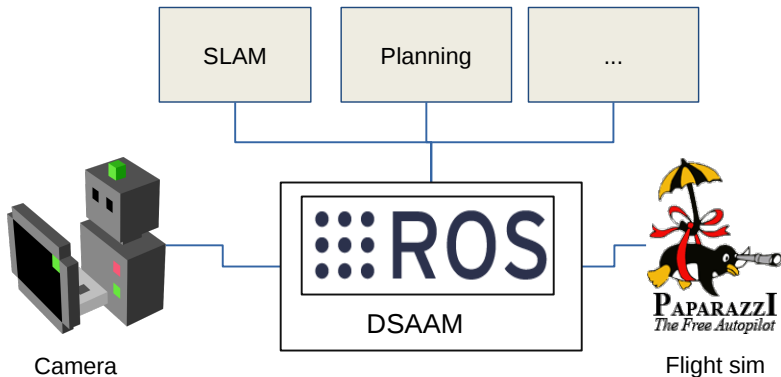
- ▶ (Only) provides time management (so far)
- ▶ Between com. middleware and app.
- ▶ May work with any kind of message passing communication middleware
- ▶ As least intrusive as possible

Minimal requirement: Messages should include timestamp

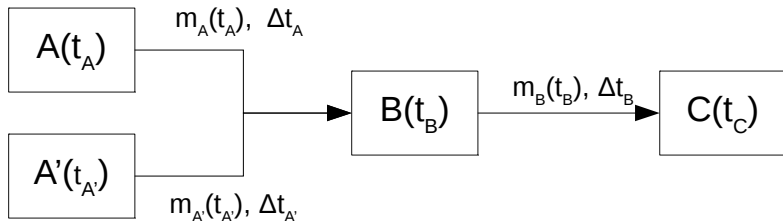
Current status:

- ▶ POC with raw threads and ROS
- ▶ native Python & C++ implementation

PRECIDRONE



INTRODUCTION



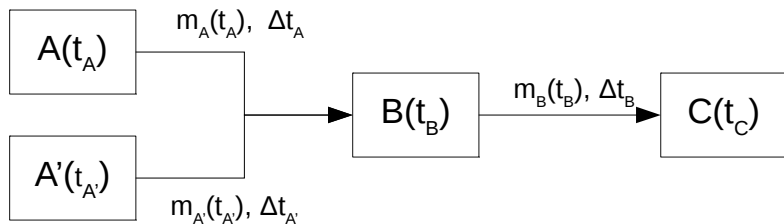
Node:

- ▶ has input and output message flows
- ▶ and it's own clock

Message flow:

- ▶ Message exchange between nodes
- ▶ Message are timestamped
- ▶ Has a message rate (evt. dynamic)

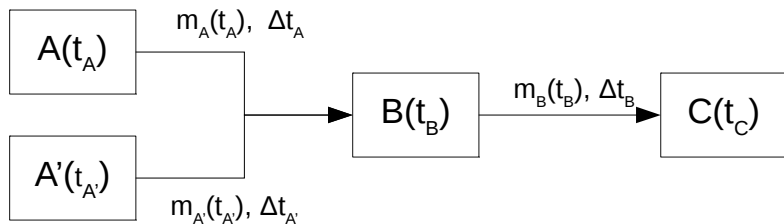
TIME MANAGEMENT (1)



Node can only advance time if all messages have been processed up to that time:

$$m_B(t), \Delta t_B \Rightarrow \text{next incoming message for B at } t' \geq t \quad (1)$$

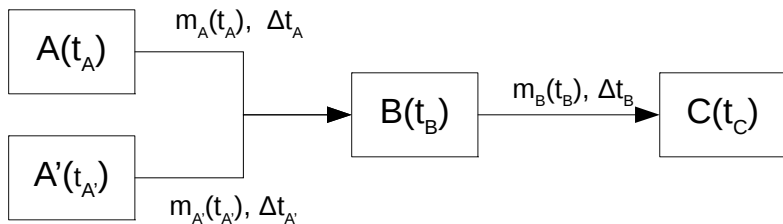
TIME MANAGEMENT (2)



Next published message timestamp can be deduced from the previous one:

$$m_B(t), \Delta t_B \Rightarrow \text{next outgoing message } m_B(t + \Delta t_B) \quad (2)$$

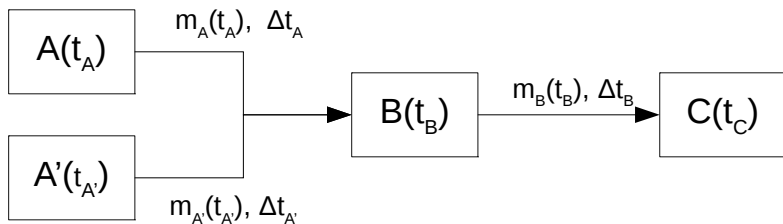
TIME MANAGEMENT (3)



Messages are processed chronologically:

$$t < t' \Rightarrow m(t) \text{ processed before } m'(t') \text{ by the node} \quad (3)$$

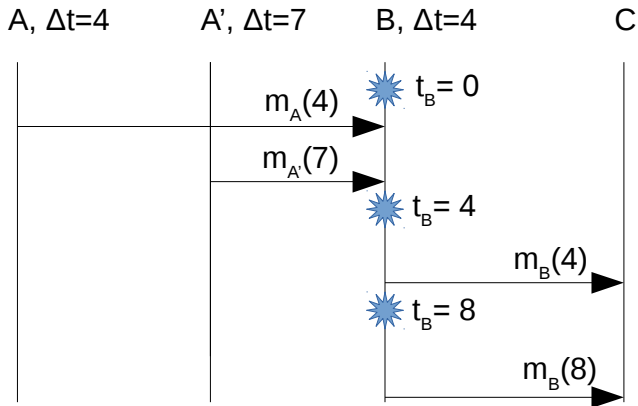
TIME MANAGEMENT (4)



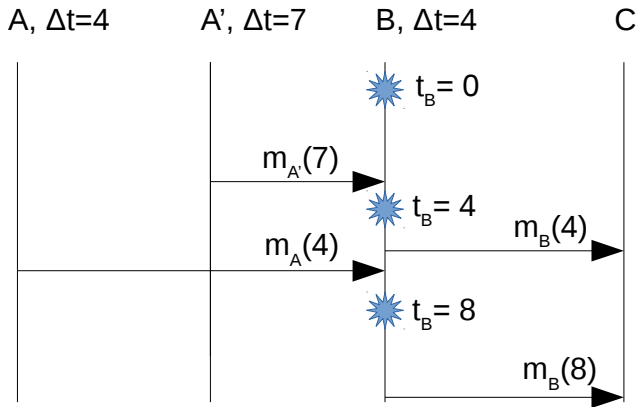
No message shall be lost:

B emits message \Rightarrow receivers queues are not full (4)

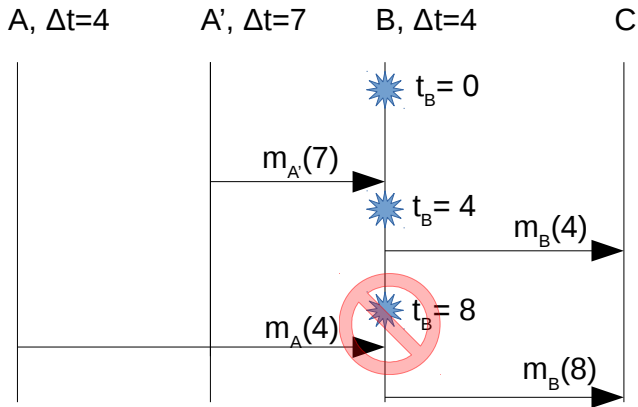
EXAMPLE (1)



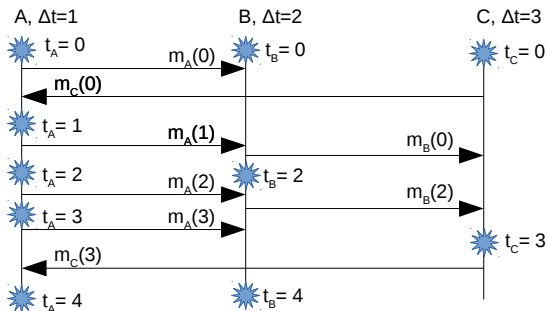
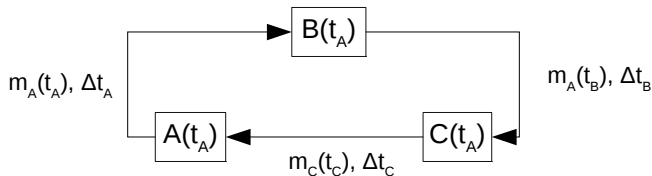
EXAMPLE (2)



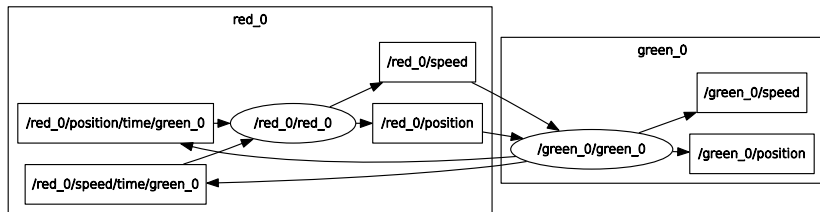
EXAMPLE (3)



EXAMPLE WITH CYCLES



ROS EXAMPLE



DUMMY EXAMPLE : ROS API IN PYTHON (1)

```
from geometry_msgs.msg import PointStamped
from dsam.ros import RosNode, Time

def process_pos(name, message, next_message_time):
    #process message

node = RosNode('B', start_time=Time(0), dt=Time(4),
              queue_size=3)

node.setup_publisher('/B/position', PointStamped,
                    dt=Time(4),
                    subscribers=['C'])

node.setup_subscriber('/A/position', PointStamped,
                     callback=process_pos,
                     dt=Time(2))
```

DUMMY EXAMPLE : ROS API IN PYTHON (2)

```
(...)
```

```
node.init_ros()
```

```
while true:
```

```
    t = node.next()
```

```
    if t > node.t + node.dt:
```

```
        m, t = simulation_step()
```

```
        node.send('/C/position', m, t)
```

SUMMARY: WHAT'S GOOD

- ▶ Decentralized time management
- ▶ Low overhead, no additional copy
- ▶ Can be implemented with any p2p message passing middleware
- ▶ Possibility of dynamic message rate
- ▶ Easy to integrate
- ▶ (Almost) transparent for apps. using "raw" middleware
- ▶ Benefits from the existing middleware ecosystem

SUMMARY: WHAT'S BAD

Difficult to fix:

- ▶ Cannot be used with mixed middlewares
- ▶ Not compatible with event based simulation (HLA is)

Could be improved or added:

- ▶ Message ACK overhead
- ▶ No simulation management
- ▶ Static flows, no dynamic reconfiguration (cf prev. point)
- ▶ No realtime/best effort mode

QUESTIONS

For the audience:

- ▶ Is repeatability a concern for you?
- ▶ Do you find the proposed solution interesting?
- ▶ If not, why?
 - ▶ Any game stoppers?
 - ▶ Missing features?
- ▶ ... vs HLA?